



Forensics for System Administrators

Advanced Memory Analysis - Dealing with
Malicious Code

Klaus Möller

WP8-T1

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Public

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Agenda - Advanced Analysis



- Creating Timelines
- User Mode Rootkits
 - User mode rootkits
 - Injected libraries
- Kernel Mode rootkits
 - Kernel modules
 - Syscall & interrupt tables
- Malware/Code extraction
 - Searching for malware signatures
 - Extracting code from kernel
 - Extracting code from processes



Timeline Creation

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Timeline Creation

Live
Demo

- Lots of timestamps in memory dumps
 - Process creation time (pslist) /termination time (psscan)
 - Thread creation time (pslist)
 - Driver/DLL/EXE compile time
 - Network socket creation time (netstat, netscan, sockets, etc.)
 - Memory resident registry key last write time (printkey et. al.)
 - Memory resident event log entry creation time
- Volatility 2 & 3: **timeliner** plug-in
 - Collects output from all plug-ins that produce timestamps and sorts them

```
> vol.py -f mem.img timeliner
```

User Mode Rootkits

User Mode Rootkit

- Rootkit - Software that hides adversary presence/activities, like
 - Processes, files, directories,
 - Network connections (sockets)
 - Free disk space, etc.
 - Interfaces in promiscuous mode (i. e. sniffers)
 - User Mode - All manipulations are in user space
 - System programs: `ps`, `ls`, etc.
 - System libraries: `libc`, `kernel32.dll`, etc.
 - Code injected into (every) process
 - Does not manipulate kernel data
 - Adversary usually still has/needs root privileges
 - With clean tools/libraries, still a good chance to detect them
- 
- Easy to find with persistent storage analysis

Manipulated System Libraries

- Rootkit functions are located in a shared library/executable
- Windows
 - Executable (`.exe`) or library (`.dll`) is injected (i.e. mapped into the address space) of the infected process
 - Overwrites OS functions/pointers or is executed in a separate thread
- Linux/Unix/MacOS
 - Setting of the `LD_PRELOAD` environment variable
 - Entries in `/etc/ld.so.preload`, `/etc/ld.so.conf`, or `/etc/ld.so.conf.d/*`
 - Direct manipulation of the `ld.so` code (the dynamic linker)
 - Malicious functions are found before regular OS functions, and linked instead of them

Searching for Libraries

- What to look for? (from Malware Analysts Cookbook)
 - Shared libraries/DLLs with suspicious names or names that haven't been seen before
 - Shared libraries/DLLs with common names that are loaded from a non-standard directory
 - For example C:\WINDOWS\sys\kernel32.dll (correct: system32)
 - Shared libraries/DLLs that allow access to protected resources or otherwise alter system security
 - Legitimate Shared libraries/DLLs that are out of context
 - I. e. libraries that are not malicious, but don't belong into certain executables, like network or security libraries in a program without network or security functionality

Example: Linux bedevil (bdvl) Rootkit

- Linux LD_PRELOAD Rootkit
 - See: <https://github.com/Error996/bdvl>
 - Hides processes with certain GID
 - Hides network connection coming from certain ports
 - Writes into /etc/ld.so.preload and patches the dynamic linker ld.so

Live
Demo



```
os151:~/bdvl # LD_PRELOAD=./build/bdvl.so.x86_64 sh -c './bdvinstall build/bdvl.so.*'

I found my sensitive side, and it has a rash.
No SELinux.
Creating installation directory.
Copied: bdvl.so.x86_64
Patching dynamic linker.
Patched ld.so: 1
Installed.

ICMP backdoor up.
PAM username: super
Accept backdoor port:
25606
Hidden port(s):
39389, 41172

Unhappy with something?
FFTUBJHVJUXHESN=1 sh -c './bdv uninstall'
os151:~/bdvl #
```

```
os151:~/bdvl # FFTUBJHVJUXHESN=1 sh -c './bdv uninstall'

Boing! It's a lightbulb! Boing!
Killing ICMP backdoor
Reverting ld.so
Eradicating directories
Uninstalling your ass
Removing preload file
Removing symlink sources
Removing other bdvl paths
Done.
```

Example: Finding bedevil

Live
Demo

- Regular libapparmor resides in /usr/lib64
- Nothing should be mapped from /lib on a 64 Bit system
- Besides, apparmor is not active on the system
- Volatility 2/3
 - Windows: `dlllist` / `windows.dlllist`
 - Linux: `linux_proc_maps`, `linux_library_list`, `linux_ldrmodules`
 - Mac OS: `mac_ldrmodules`



Task	Pid	Load Address	Path
dbus-daemon	792	0x00007ffd42585000	linux-vdso.so.1
dbus-daemon	792	0x00007f2410986000	/usr/lib64/libdbus-1.so.3
...			
dbus-daemon	94779	0x00007f1d6db2e000	/lib/apparmor/libapparmor.so.x86_64
...			

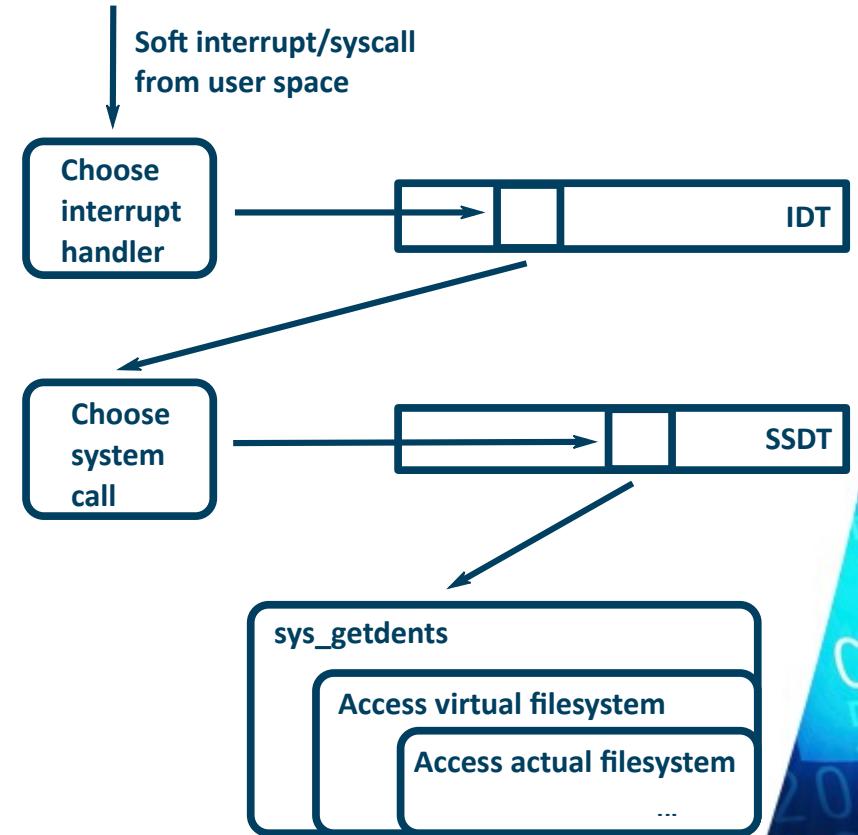
Kernel Mode Rootkits

Kernel Mode Rootkits

- Same functionality as user space rootkits
- But some/all code is now executed in kernel space
 - System calls return incomplete or false data
 - No way to detect kernel rootkits with clean tools/libraries
 - Not as long as the compromised kernel is running
- Memory analysis does not rely on the kernel
 - Instead, we examine the in-memory data structures directly
- Two common (among other) approaches of rootkit authors
 - Manipulating the system call or interrupt table
 - Loading a malicious kernel module/driver

System Call & Interrupt Table

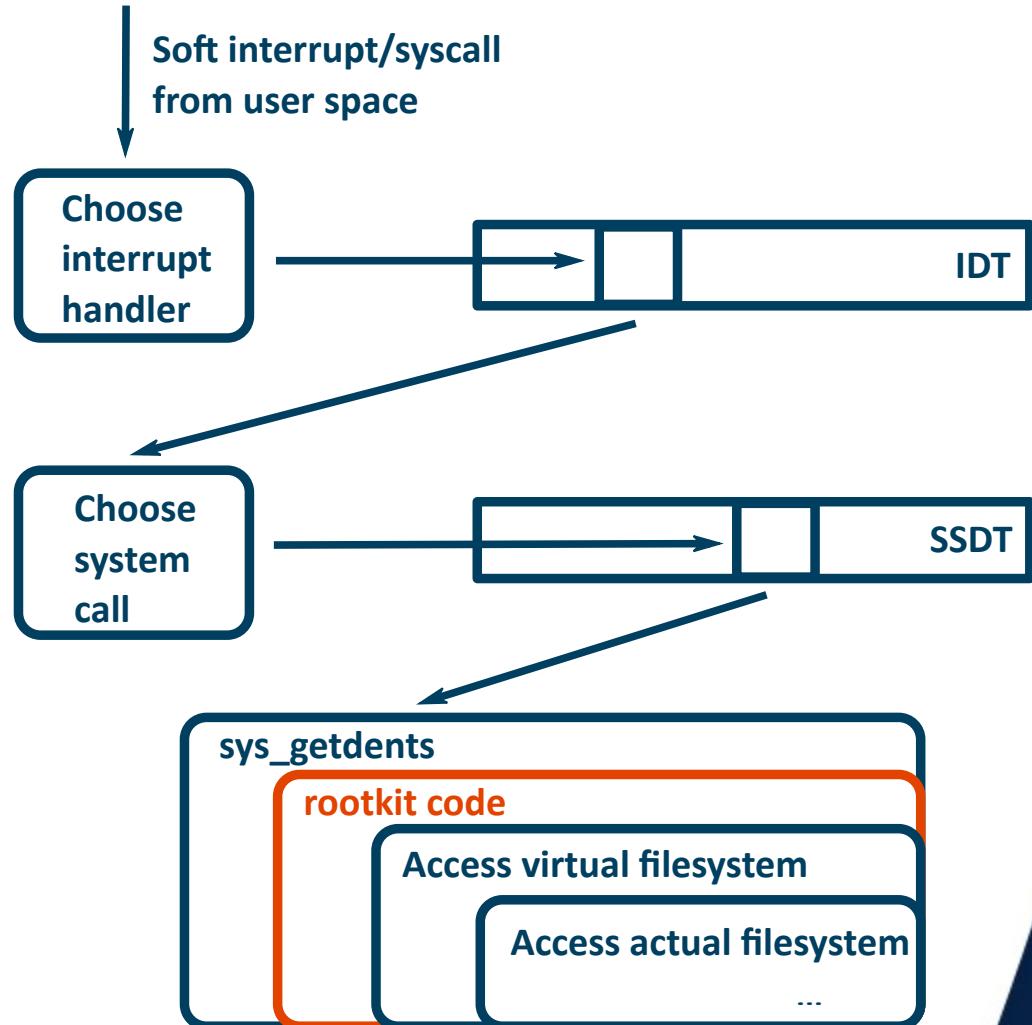
- Interrupt descriptor table (IDT) (Intel)
 - Table indexed by processor interrupt number
 - Entries are basically pointers to functions that are called when the interrupt is raised
- Service Descriptor (or System Call) Table (SSDT) (Windows/Linux)
 - Program code can raise interrupts, usually through special instruction
 - Called *soft interrupts* to differentiate from hardware generated ones (disks, keyboard, mouse, etc.)
 - Supplied with this interrupt is a number, used to look up in the SSDT
 - SSDT is indexed by this *system call number*
 - Entries are pointers to the system call functions
- Rootkits manipulate this lookup process to hide



Source: Andreas Bunten: *Unix und Linux Kernel-based Rootkits* (2004)

Direct manipulation of the system call code

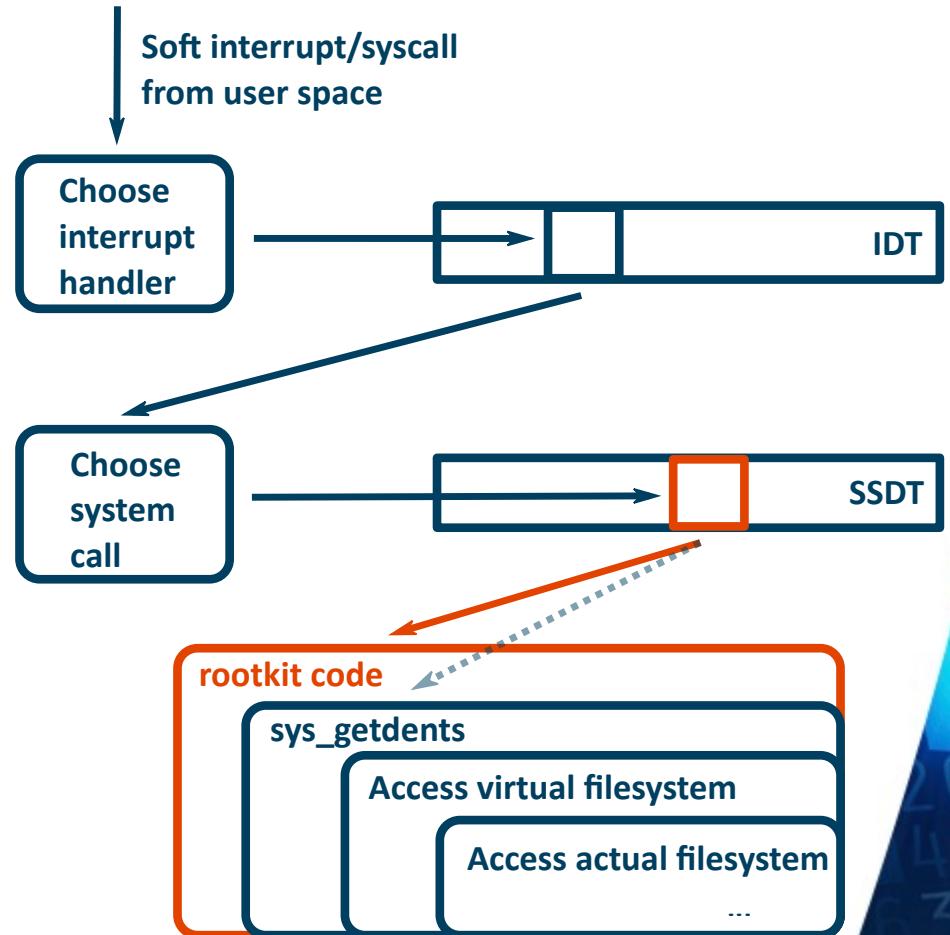
- IDT and SSDT entries and base addresses unchanged
- Syscall/Interrupt handler code unchanged
- Rootkit code is inserted into the actual syscall function



Source: Andreas Bunten: *Unix und Linux Kernel-based Rootkits* (2004)

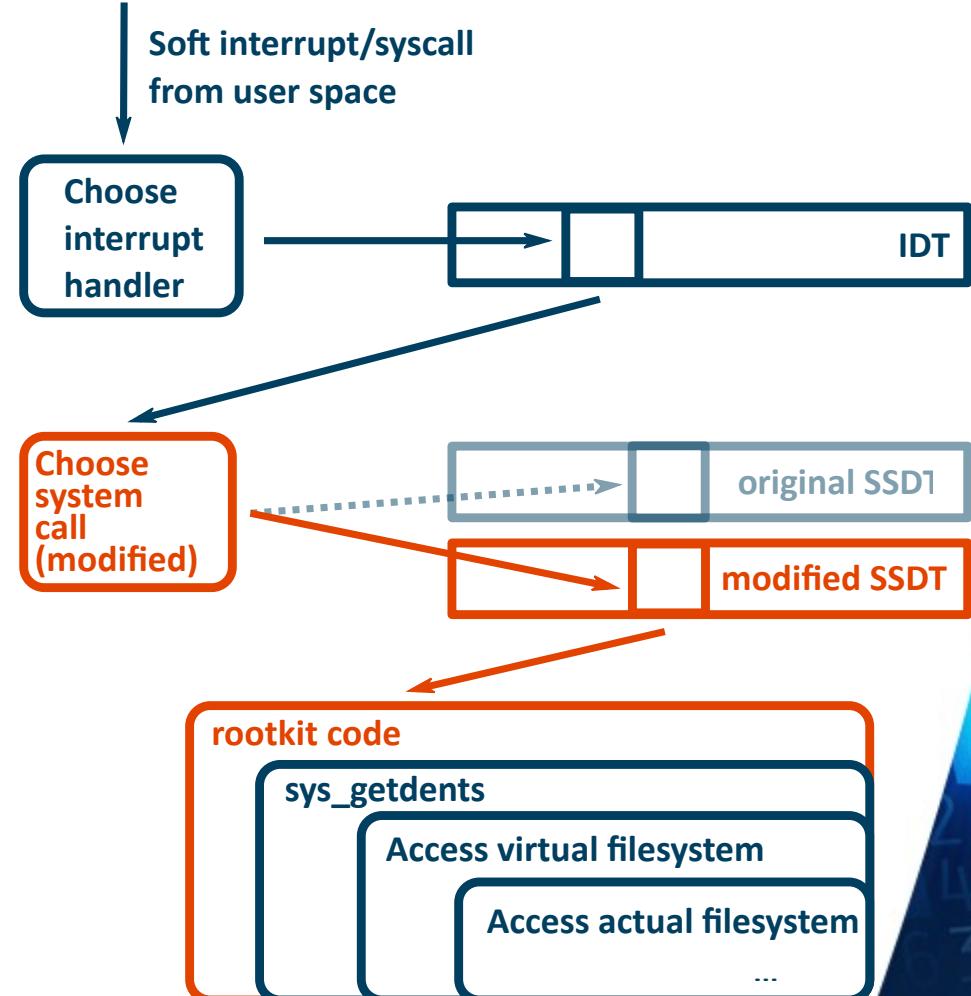
Replace SSDT entry

- Replace SSDT entry with a pointer to a function from the rootkit
 - May in turn call the original system call function, rootkit works than as a wrapper around the original system call code
 - Entry may stand out, when it points into another memory region
- IDT and soft interrupt handler unchanged
- Base SSDT address unchanged



Modify Soft Interrupt Handler

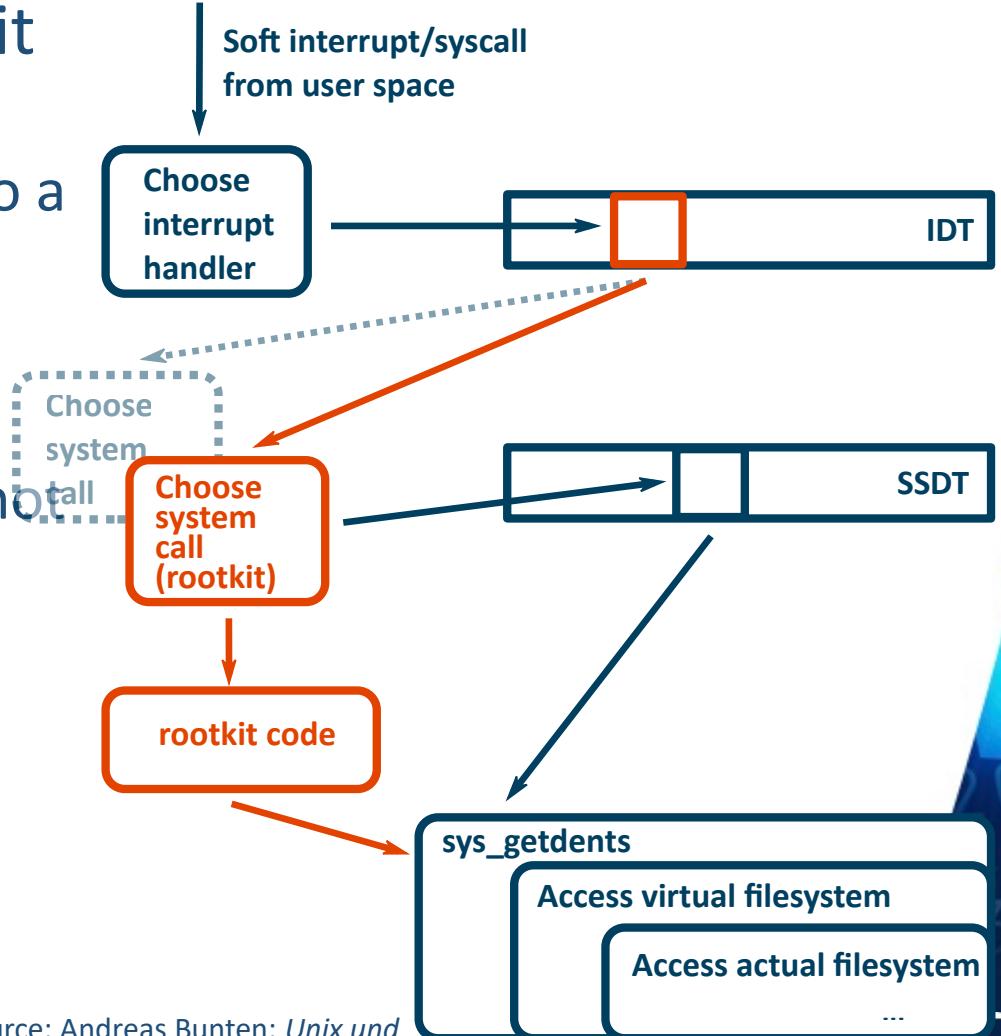
- Rootkit interrupt handler uses its own SSDT with manipulated functions
- Entries in the original SSDT are looking homogenous - in fact the SSDT itself is not changed
- Pointers from the rootkits SSDT still point to rootkit wrappers
- Search has to inspect the soft interrupt handler code to find rootkit
- Or to find the modified copy of the SSDT and compare to original



Source: Andreas Bunten: *Unix und Linux Kernel-based Rootkits* (2004)

Replace IDT Soft Interrupt Handler Entry

- Soft interrupt entry points to rootkit handler
 - Usually stands out as it may point into a different memory region than other interrupt handlers
- Rootkit handler uses original SSDT
 - Thus, SSDT and syscall functions are not manipulated (and look genuine)
- But bypasses SSDT for hooked functions
 - These rootkit functions may call the functions from the unmodified SSDT



Source: Andreas Bunten: *Unix und Linux Kernel-based Rootkits* (2004)

Volatility IDT/SSDT Plugins

- Plugins to inspect the IDT and SSDT (or syscall table)
 - Windows: `idt`, `ssdt` / `windows.ssdt`
 - Linux: `linux_idt`, `linux_check_evt_arm`, `linux_syscall`,
`linux_syscall_arm` / `linux.check_idt` `linux.check_syscall`
 - Mac OS: `mac_idt`, `mac_syscalls`, `mac_check_syscall_shadow` /
`mac.check_syscall`
- Critical functions (page faults, etc.) should always be handled by the core OS or standard modules
- Lots of more function pointer tables in the OS that can be used by rootkits
 - And more plugins to find them

Kernel Modules & Rootkits

- Kernel rootkit functions often bundled into a kernel module
 - Installation uses standard module loading mechanisms of the OS
- Often user space process to communicate with the kernel parts
 - Socket/handle between process and kernel module
 - User space process can be detected (sometimes not hidden)
- To list loaded kernel modules or drivers (Volatility 2/3)
 - Linux: `linux_lsmod` / `linux.lsmod`
 - Windows: `modules`, `modscan`, `driverscan`, `driverirp` / `windows.modules`, `windows.modscan`, `windows.driverscan`, `windows.driverirp`
 - Mac OS: `mac_lsmod` / `mac.lsmod`

Example: Windows Laqma rootkit

Live
Demo

- Rootkit installs **lanmandrv.sys**
 - Path suspicious: \windows\system32 instead of \windows\system32\drivers
 - Genuine LANMAN driver is called lanman.drv
 - Towards the end/beginning of the list, may be a sign of later loading of the driver/module

```
noob@v2os153:~/volatility> ./vol.py -f img/laqma.vmem modules
Volatility Foundation Volatility Framework 2.6.1
Offset(V) Name Base Size File
-----
0x810dbe68 ntoskrnl.exe 0x804d7000 0x1f6280 \WINDOWS\system32\ntkrnlpa.exe
0x810dbe00 hal.dll 0x806ce000 0x20380 \WINDOWS\system32\hal.dll
0x810dbd98 kdcom.dll 0xfc99b000 0x2000 \WINDOWS\system32\KDCOM.DLL
0x810dbd28 BOOTVID.dll 0xfc8ab000 0x3000 \WINDOWS\system32\BOOTVID.dll
0x810dbcc0 ACPI.sys 0xfc36c000 0x2e000 ACPI.sys
...
0xff381bd8 sysaudio.sys 0xf337d000 0xf000 \SystemRoot\system32\drivers\sysaudio.sys
0xff2837e8 Fastfat.SYS 0xf2ef5000 0x23000 \SystemRoot\System32\Drivers\Fastfat.SYS
0x80fae8b0 lanmandrv.sys 0xfc29000 0x2000 \??\C:\WINDOWS\System32\lanmandrv.sys
0xff147970 kmixer.sys 0xf2ecb000 0x2a000 \SystemRoot\system32\drivers\kmixer.sys
noob@v2os153:~/volatility>
```

Malware

Searching for Malware Signatures

- VAD (Virtual Address Descriptor) are descriptor structures for memory regions in Windows
- **Malfind** plug-in searches the VAD blocks for code that may have been injected by malware
 - An **MZ** header in a block is almost always suspicious as it is normally not included
 - Maybe a whole executable was injected en-block
 - Header list can be limited to “suspicious” ones with -W
 - I. e. those containing “MZ”
- Suspicious blocks can be extracted with -dump-dir= (or -D)
- Files can than be reverse engineered or scanned for malicious code
 - I. e. AV-scanner or virustotal.com

Example: Zeus Banking Trojan (aka Zbot)

Live
Demo

- Scanning (and dumping) the virtual memory image

```
> ./vol.py -f img/zeus.vmem malfind -W -D dumpdir
> cd dumpdir; ls -l
noob@v2os153:~/volatility/dumpdir> ls -l
insgesamt 3648
-rw-r--r-- 1 noob users 155648 9. Mai 20:34 process.0x80f60da0.0x1000000.dmp
-rw-r--r-- 1 noob users 155648 9. Mai 20:34 process.0x80f94588.0x12d0000.dmp
...
> sha1sum process.0x80f60da0.0x1000000.dmp
a2c13b59832d1cdb25ba2f6ba8f6cacd044c51a8 process.0x80f60da0.0x1000000.dmp
```

Example: Zeus Banking Trojan (virustotal upload)

- Surprise: It's Zbot;)

Live
Demo

The screenshot shows a virustotal analysis page for a file named 'process.0x80f60da0.0x1000000.dmp'. The file has a size of 152.00 KB and was uploaded 2 months ago. The analysis summary indicates that 57 security vendors flagged the file as malicious, while 69 flagged it as safe. The file is identified as an EXE file.

Community Score: 57 / 69

File Details:
8d57c9e914214a45ca913544c1d9a09312df5c4b8733049442176f1e26915990
process.0x80f60da0.0x1000000.dmp
invalid-rich-pe-checksum overlay peexe spreader

Analysis Summary:
57 security vendors and no sandboxes flagged this file as malicious

Detection: 57/69

Details: 69/69

Behavior: 69/69

Community: 69/69

Security Vendors' Analysis:

Vendor	Result	Engine	Description
Acronis (Static ML)	! Suspicious	Ad-Aware	! Gen:Variant.Razy.447136
AhnLab-V3	! Worm/Win32.IRCBot.C136977	Alibaba	! TrojanPSW:Win32/ShellCode.db9a6d29
ALYac	Gen:Variant.Razy.447136	Antiy-AVL	! Trojan/Generic.ASMalwS.33B24E6
Arcabit	! Trojan.Razy.D6D2A0	Avast	! Sf:Crypt-BT [Trj]
AVG	! Sf:Crypt-BT [Trj]	Avira (no cloud)	! TR/Patched.Ren.Gen
BitDefender	! Gen:Variant.Razy.447136	BitDefenderTheta	! Gen>NN.Zexaf.34232.jqZ@aqTeVHc
Bkav Pro	! W32.AIDetect.malware1	ClamAV	! Win.Spyware.Zbot-9841872-0

Code Extraction from Kernel

Live
Demo

- Volatility plug-ins to dump kernel modules (V2 only)
 - Windows: `moddump`
 - Linux: `linux_moddump`
 - Mac OS: `mac_moddump`
- Use base address from `modules` or `*_lsmod` plug-ins for `-b` option
 - Without, all modules will be dumped

```
> ./vol.py -f img/laqma.vmem moddump -b 0xfca29000 -D .
Volatility Foundation Volatility Framework 2.6.1
Module Base Module Name          Result
-----
0x0fca29000 lanmandrv.sys      OK: driver.fca29000.sys
> sha1sum driver.fca29000.sys
01e53955f889fb317cd4981d03c622d8b60dc2f5  driver.fca29000.sys
```

Code Extraction from Processes

Live
Demo

- Plugins to dump process memory to file(V2 only)
 - Windows: **procdump**
 - Linux: **linux_procdump**
 - Mac OS: **mac_procdump**

```
> ./vol.py -f img/laqma.vmem pslist | grep lanm
Volatility Foundation Volatility Framework 2.6.1
0xff3825f8 lanmanwrk.exe          1180    1060      2      75      0      0
2010-08-15 19:09:12 UTC+0000
> ./vol.py -f img/laqma.vmem procdump -p 1180 -D procdumpdir/
Volatility Foundation Volatility Framework 2.6.1
Process(V) ImageBase Name           Result
-----
0xff3825f8 0x00400000 lanmanwrk.exe      OK: executable.1180.exe
noob@v2os153:~/volatility> sha1sum procdumpdir/executable.1180.exe
d84f38096f040b167919bb8a639837305bf690bf  procdumpdir/executable.1180.exe
```

Wrapping Up

What's Next?

- We've seen only the very beginning
 - Lots of more plug-ins for lots of more ways for adversaries to hide
 - Read the volatility blog, documentation, books, etc.
- Practice, practice, practice
 - Try analysing some of the sample images with real malware, or do some memory acquisition from your own systems
 - See references section
 - For production systems, ask for permission first!
- Don't forget: There are more ways/tools for forensic analysis
 - Persistent storage analysis (Autopsy, next webinars)
 - Network data analysis (Wireshark)
 - Reverse engineering (Ghidra, Radare 2)
 - And more, ... even very simple ones (strings, objdump, ...)



Finally

- The situation is hopeless - but not serious;)
- For every attack, there's a defence (and vice versa:(
- There will always be traces the adversary forgot, couldn't avoid or did not care about

Seek and ye shall find!

Thank you

Any questions?

Next Webinar: ***Persistent Storage Forensics I***

May 25th, 2022

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