



# Another Nail in the Coffin: OpenPGP Key Servers Gone Rogue

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# Outline

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## Introduction

- Use Cases for Key Servers
- SKS – Synchronising Key Servers
- OpenPGP basic structure + RSA basics
- **Why we need revocation**

## Problems

- Missing Integrity Check
- Equivocation
- **The current concept and it's implementation are broken**

## Solution

- Certificate Transparency Log
- PGP Revocation Bytes
- **How to design a Revocation Service**

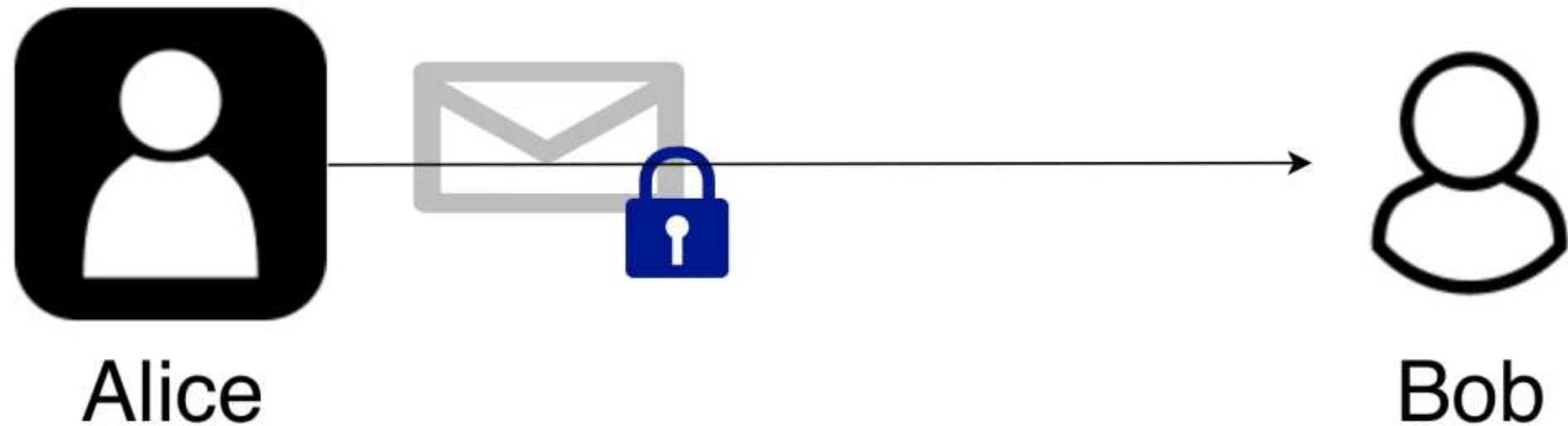
## Intro

- Why do we need to exchange keys?
- What to I do when my keys got stolen?
  - Generate a Revocation Certificate beforehand
- How do PGP packets work?

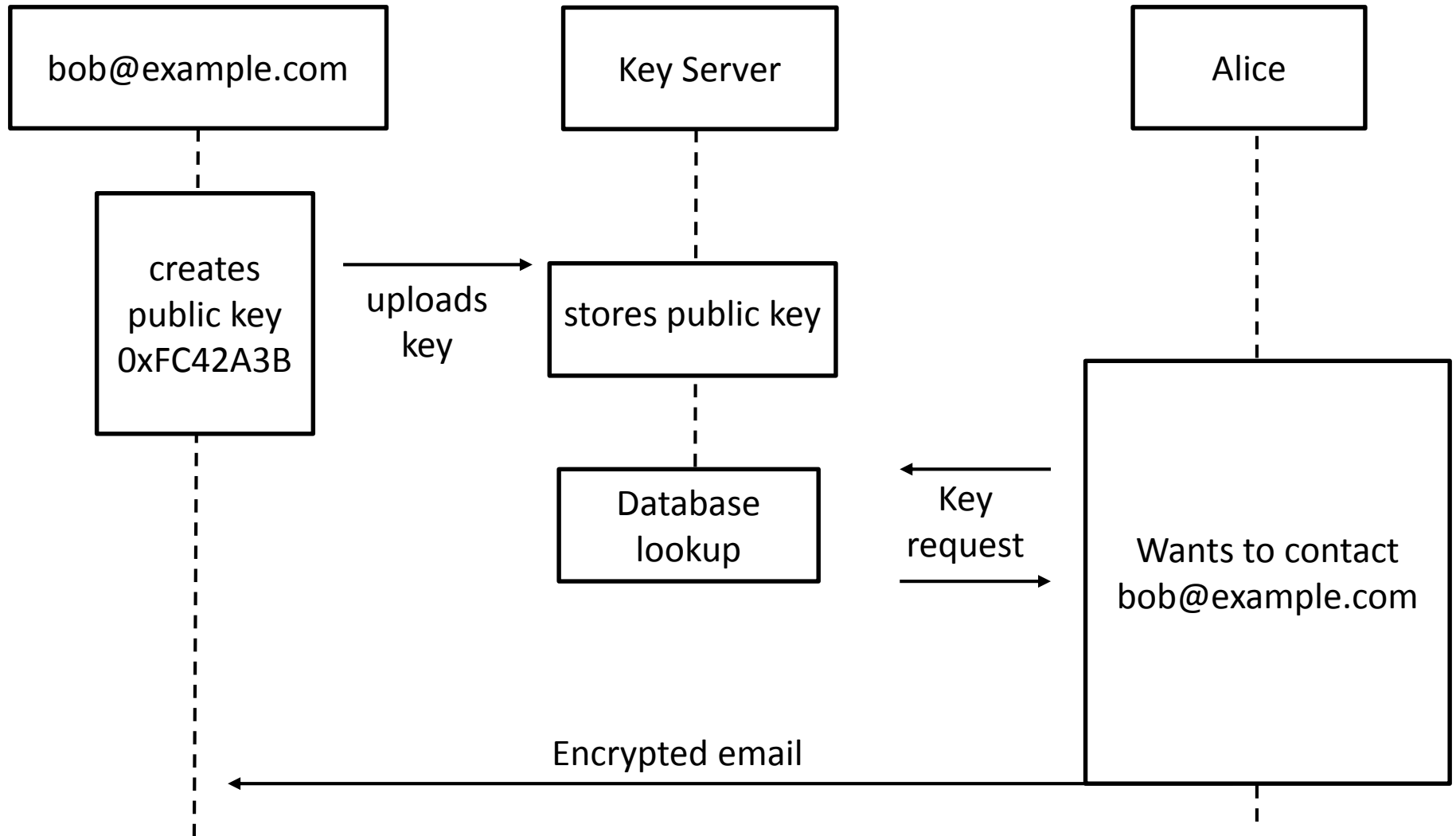


## Sending and encrypted Email

- Alice needs a key for sending an encrypted Email to Bob



# Promise: Availability without prior contact



## Key Servers' promise: Availability of Cryptographic Keys

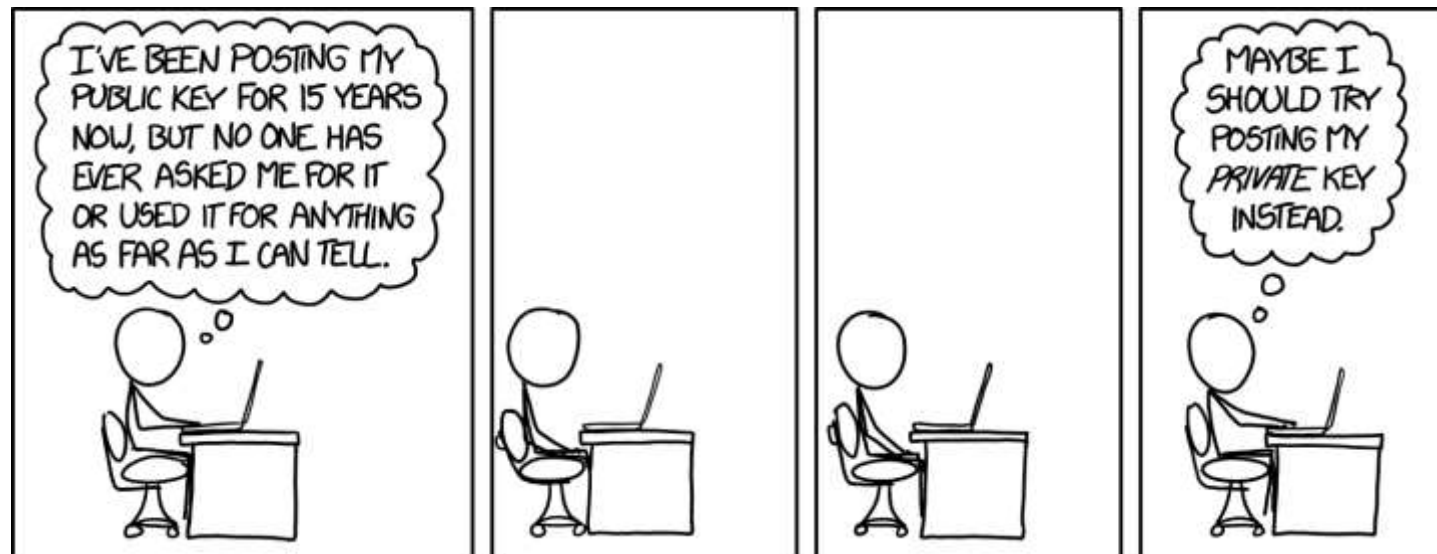
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- Receiving the actual key material by UID or fingerprint
- Mitigate Deletion of Keys on one Key Server by Re-Distribution
- Update certificate, e.g. new sub-keys, new UIDs, new certifications
- **Check for revocations**

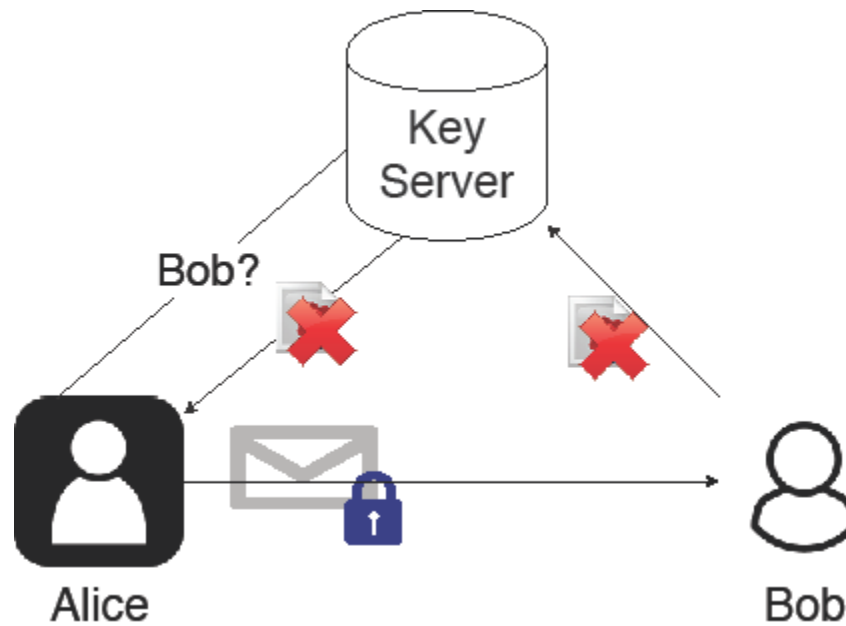
**You have to check for integrity yourself**

## What if my **private key** gets compromised?

- When an attacker could successfully obtain the private key of Bob, Bob needs to **revoke** his key.  
Notice - Bob should generate a **revocation certificate for his key in advance**

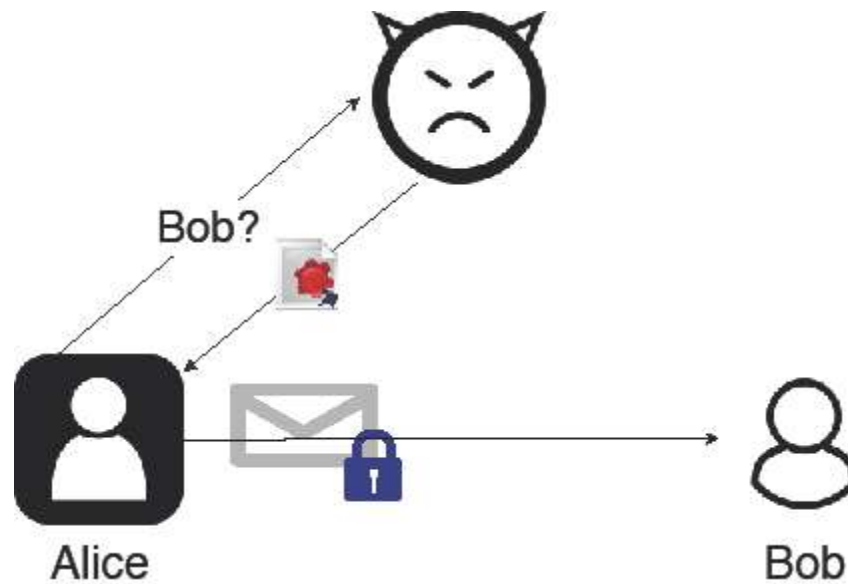


## Revocation – The happy path

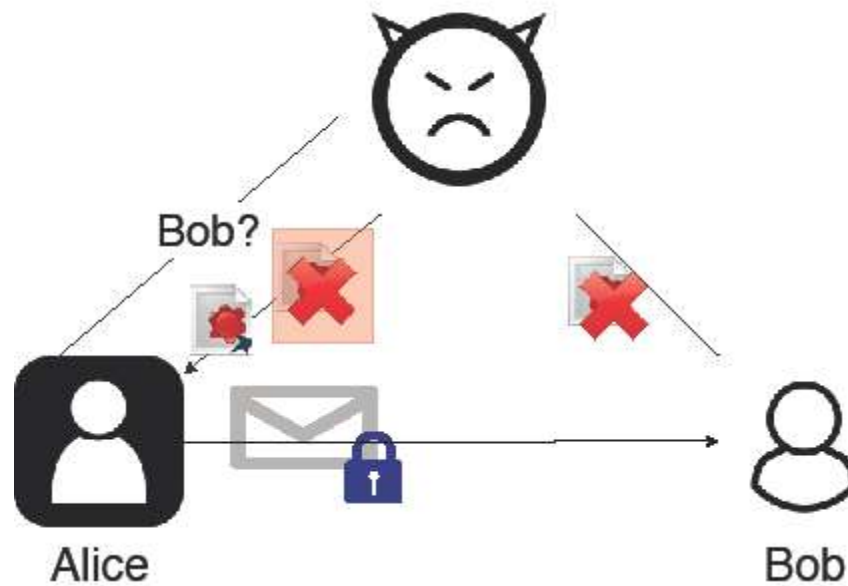




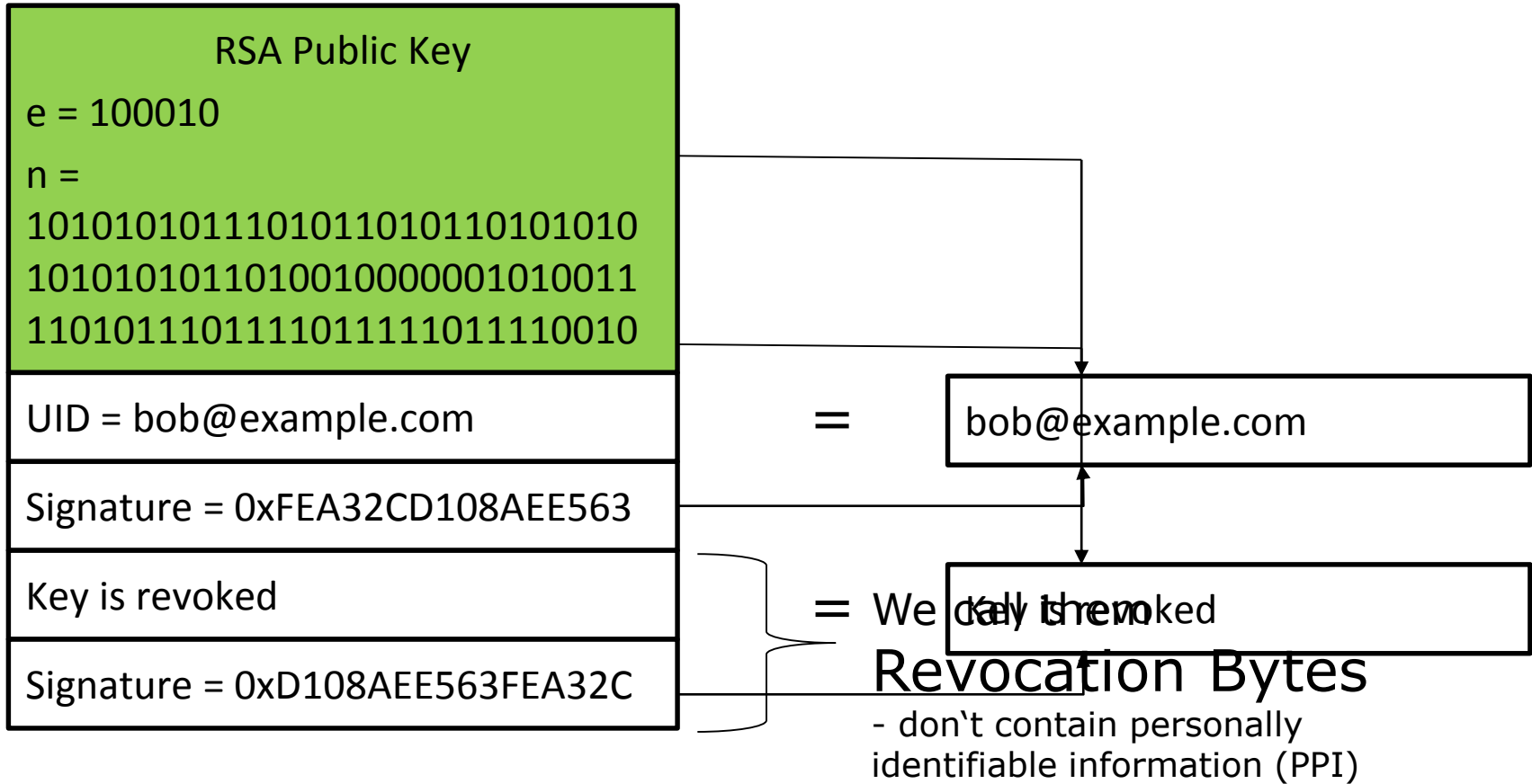
## Revocation – The evil path



## Revocation – The evil path, Equivocation



# Validation



# Invalid packets get discarded

<b>RSA Public Key</b> e = 100010 n = 1010101011101011010110101010 1010101011010010000001010011 1101011101111011111011110010
UID = bob@example.com
Signature = 0xFE32CD108AEE563
<del>Attacker says, key is revoked</del>
<del>Signature = 0xFFFFFFFFFFFFFFFF</del>

≠

Ä³korkrÄ°ss@cafÄ©p.eÄ±o

Invalid Revokation Bytes

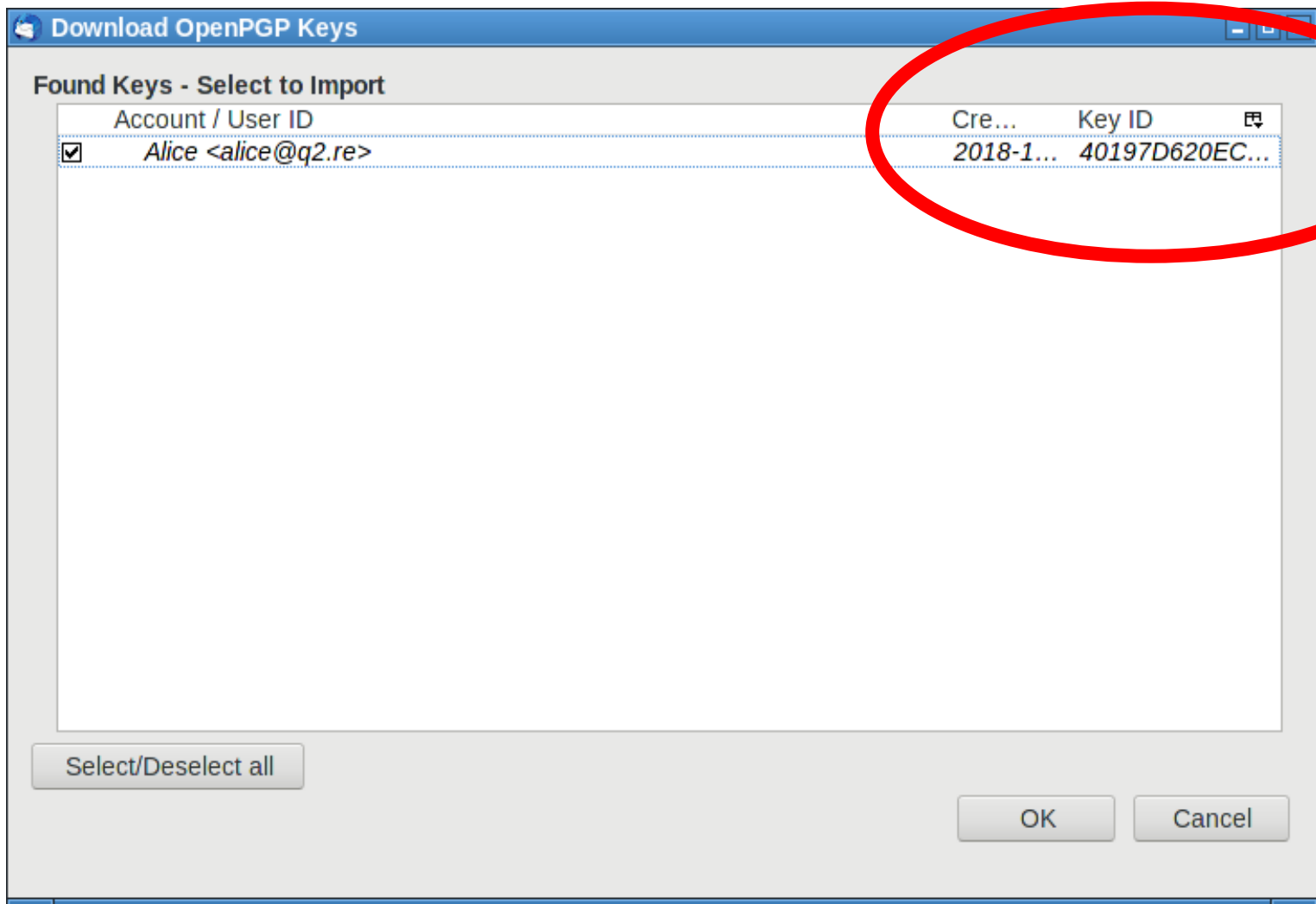
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..., so what does the key server check then?

# Anatomy of an OpenPGP Certificate

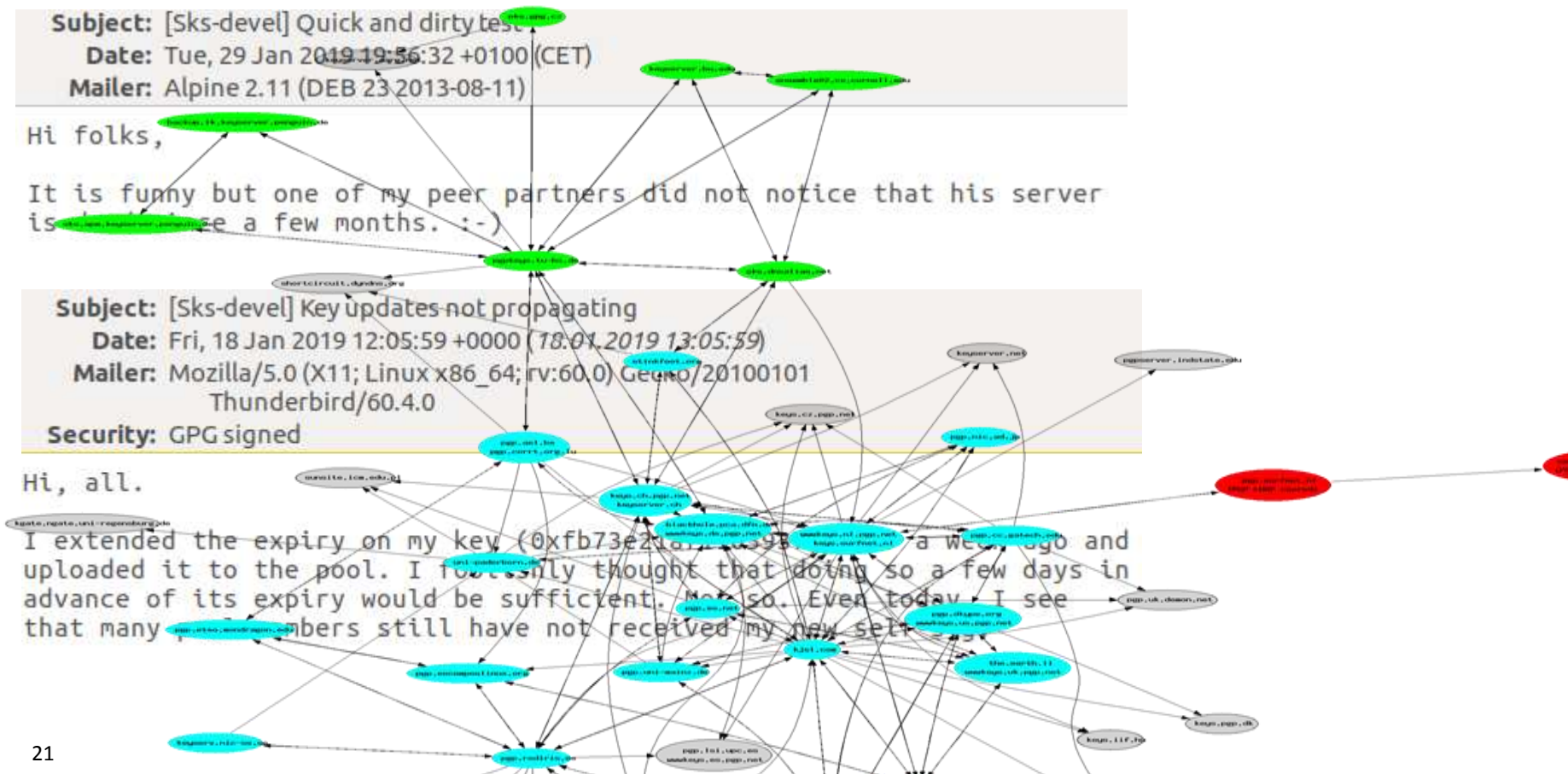
0000	88 78																			frame
0002			04																	version
0003				20																<u>sigtype</u>
0004					16															pk_algo
0005						08														hash_algo
0006							00 20													hashed_area_len
0008								16 21 04 cd 6f 2d 93 e8												hashed_area
0010	87 0d 61 17 6a d5 c7 89							b4 a0 07 47 a2 88 70 05												
0020	02 5b dd 82 a9 02 1d 00																			
0028								00 0a												unhashed_area_len
002a									09 10 89 b4 a0 07											unhashed_area
0030	47 a2 88 70																			
0034					06															hash_prefix1
0035						f9														hash_prefix2
0036							01 00													mpi_len
0038								97 61 89 af 42 3d f5 e2												eddsa_signature_r
0040	eb 95 9a 50 d2 c4 20 ce							fc 2a f7 f9 1b 72 27 33												
0050	0e 5a 4b 7a 2f 27 73 2a																			
0058								01 00												mpi_len
005a									9d 3b 4a 71 c2 2f											eddsa_signature_s
0060	1e 2b 65 f4 24 20 11 3d							45 29 ee 16 fc 61 ef 3f												
0070	fd 98 16 f7 98 e1 33 48							1e 07												

Enigmail showing a revoked certificate  
(but it has not been revoked!)



# SKS – Synchronising Keyserver

- Main implementation of HKP enabled daemon
- Set-reconciliation with near optimal communication
- Loosely connected set of servers in the pool





# SKS – Synchronising Keyserver – inherent problems

- Cannot delete keys
  - GDPR concerns
  - Poison keys

the pool of SKS keyserver: as anyone can upload anybody's key, and it does not allow to delete keys, it's IMHO by not compatible with GDPR.

- Equivocation
  - Strip packets

**Subject:** Re: [Sks-devel] Unusual traffic for key 0x69D2EAD9 and 0xB33B4659  
**Date:** Mon, 04 Feb 2019 10:49:30 +0100  
**Mailer:** Roundcube Webmail/1.2.3

Hi,

Don't get me wrong, but within three days I've got 450G traffic which can be assigned to sks by 99.9%. Estimated to 30 days this means 4.5T (which is in good agreement of your 2+T/Key for these two poison keys).

# SKS – Synchronising Keyserver – operational problems

## ■ Single threaded Ocaml

- Arcane architecture for high performance pools

- Maintenance

- Pool size

## ■ No integrity validation

- No semantic checks, i.e. identity

- No cryptographic checks, e.g. Invalid signatures

- Only rudimentary validation against the OpenPGP spec

```
> And how do we tell sks to use more than ONE instance for the SAME  
> database?!
```

```
You don't. SKS just isn't built that way. To get concurrency, you need  
to run multiple separate instances of SKS and configure them to gossip  
between each other. Then you can put a load balancing reverse proxy in  
front to simulate a multi-threaded server. Kristian and a few others  
have been operating this way for a while now.
```

```
> Any chance that sks will be fixed some day?
```

```
Short answer, no. SKS is effectively running as end-of-life software at  
this point. It gets emergency bugfixes but little else. The improvements  
you are talking about would require an enormous refactoring of the  
codebase, likely requiring migration to a different database engine.  
Given that there are fundamental design flaws (poison keys) that aren't  
getting fixed, performance issues just aren't on the radar. Sorry.
```

```
1. Future updates for the key will be denied, including  
legitimate ones by key holder (FreePBX team). 2. DoS is still  
possible just by accessing/fetching the key. To fix that, you'll  
have to remove the DoS packets (large user packets with random  
gibberish, not valid per OpenPGP packet spec, does not validate  
cryptographically) or the whole key. 3. Anyone can create another  
poison key at any time and there's no way to fix that without  
breaking compat, it's a fundamental flaw :-)
```

# The final breath of SKS

(\*2002 - †2019)

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I have not been running a keyserver myself since the first poison-key incident. Given the ongoing poison-key problems I am unlikely now to ever run one again, unless the entire codebase is overhauled. I have neither the time, energy nor skillset to perform this task, and it is becoming clear that nobody else does either.

This is not anyone's fault. It is just an unfortunate reality that we have to accept. SKS is end-of-life.

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What can we do?

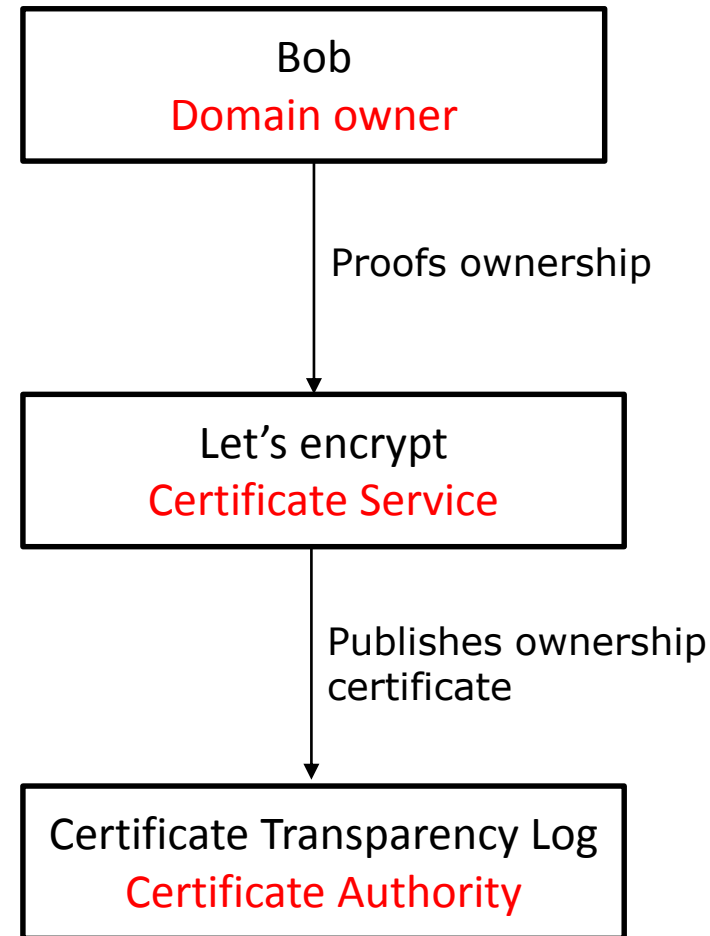
## Requirements for a Revocation Service

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- **Integrity**
  - A key server must not be able to alter data it has received.
- **Non-Repudiation**
  - The key server must be held accountable for serving manipulated data.
- **Privacy**
  - The key server must not learn who a client wishes to communicate with.
  - The uploaded data must not contain Personally Identifiable Data (PII).
- **Offline-Capability**
  - The store and forward architecture of the email ecosystem does not require the recipients to be online at the same time.
- **Timeliness**
  - Once a packet has been uploaded to a key server, it must serve it to clients requesting it.
- **Equivocation-resistance**
  - A key server must provide the same answer regardless of the identity of the client.

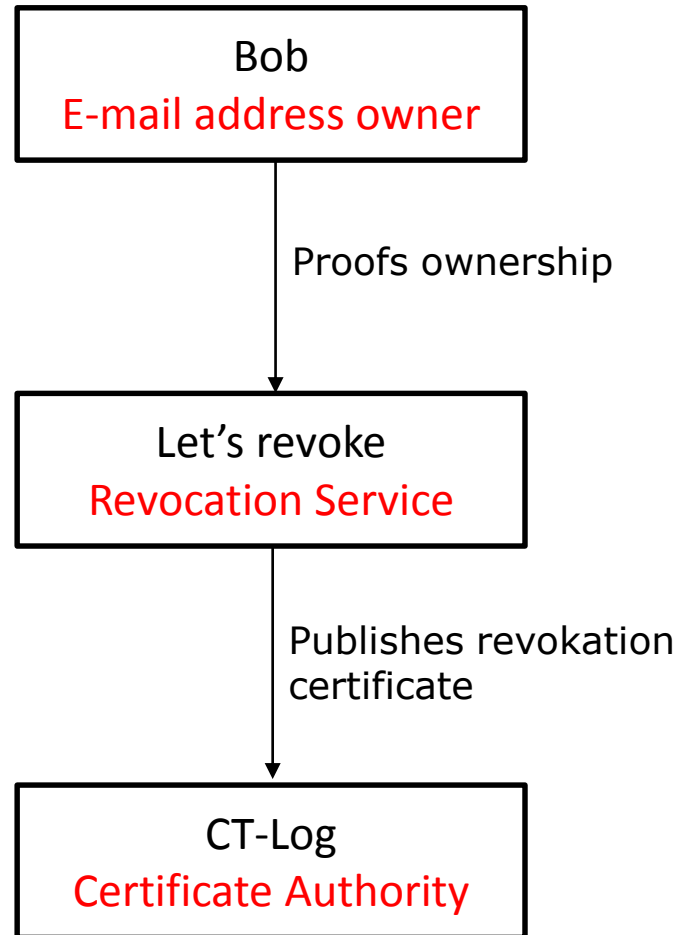
## Let's encrypt?

- A public service
- Fully automated
- Objective: Hands-out certificates for your domain, e.g. bobswebsite.de
- Runs as a script on your server
- You have to provide proof of access
  - e.g. by putting downloadable files somewhere
- Then you get a certificate for you domain
- And all major browsers accept it
- It is possible to **verify** automatically **before** publishing!
- Why not try this for emails?



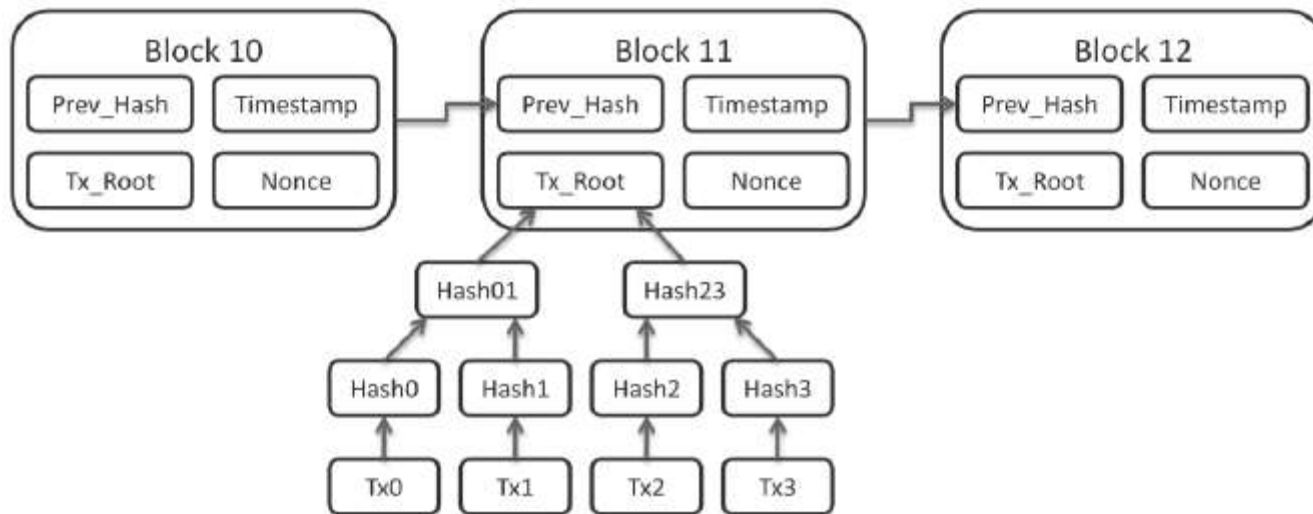
What do we need?

Let's  
revoke!



# Towards an Equivocation resistant key revocation scheme

- Publicly Verifiable Append-only Data Structure
- Expensive to operate



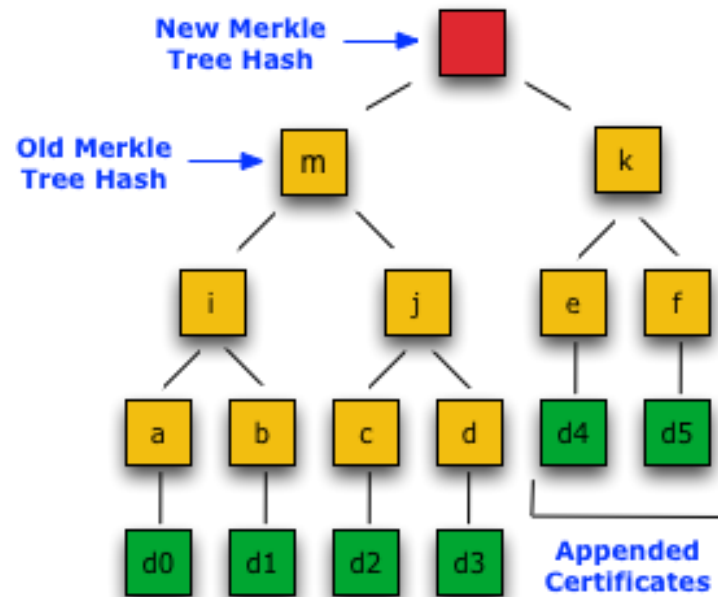


## CT-log-based Revocation Protocols

- WebPKI does not have an effective revocation scheme
- WebPKI uses Certificate Transparency to uncover wrongly issued certificates
- Google, Cloudflare, and others operate CT logs

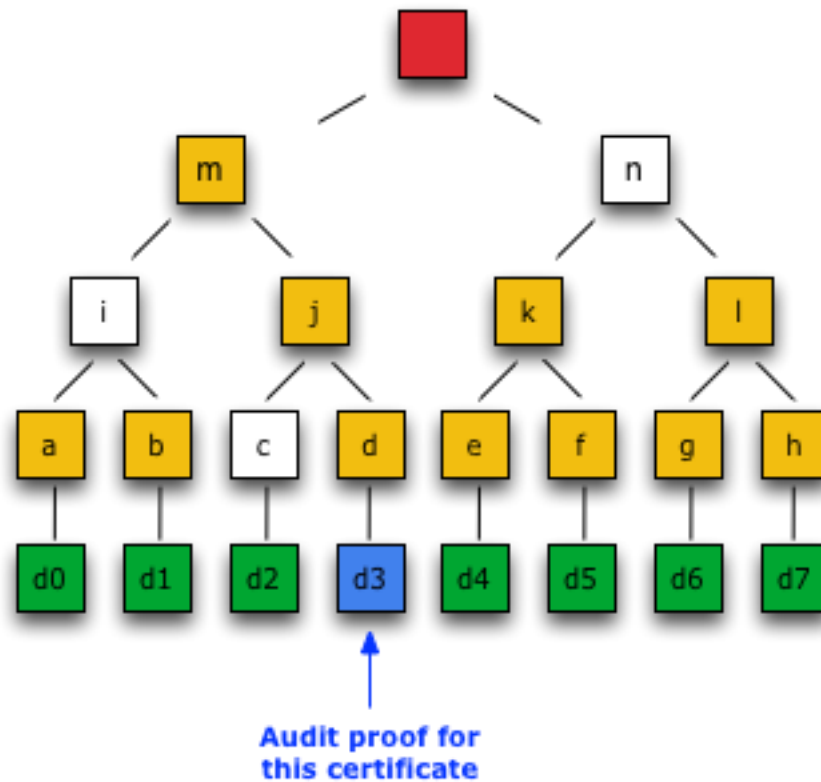


# Certificate Transparency - Merkle Tree



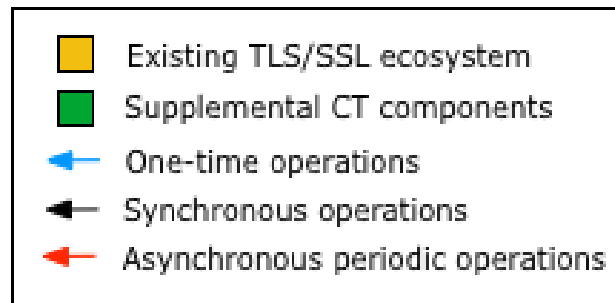
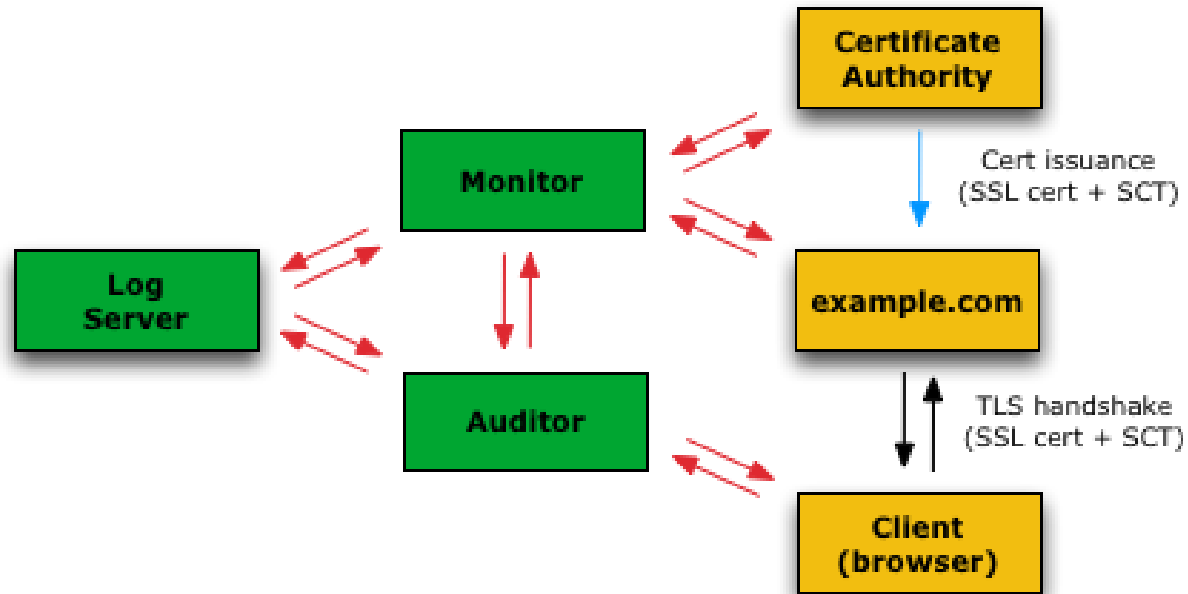
**Figure 2**

# Certificate Transparency – Audit Proof



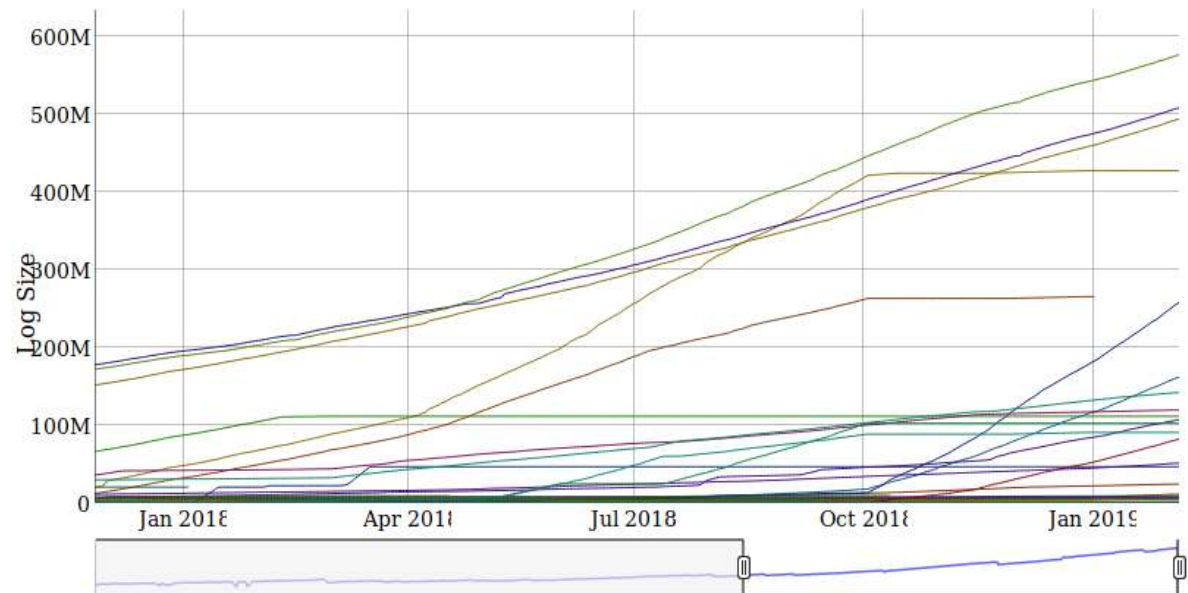
**Figure 5**

# Certificate Transparency - Infrastructure

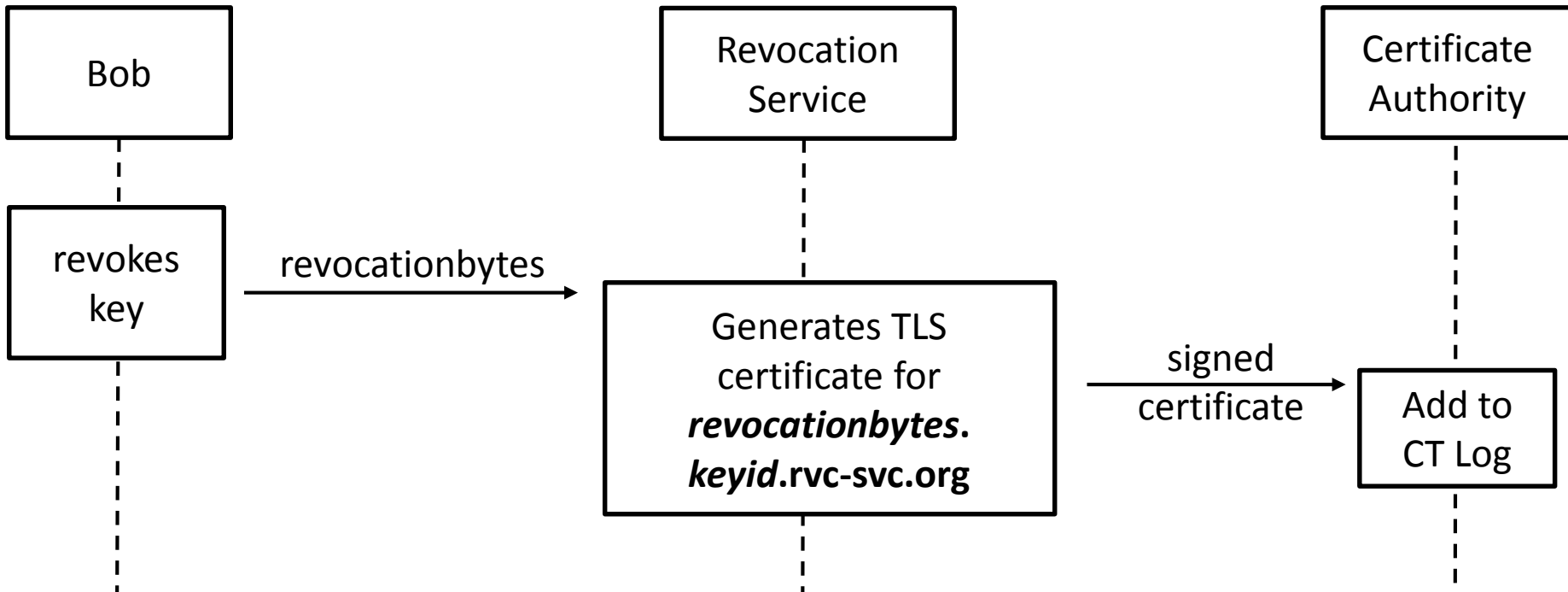


# Revocation Service

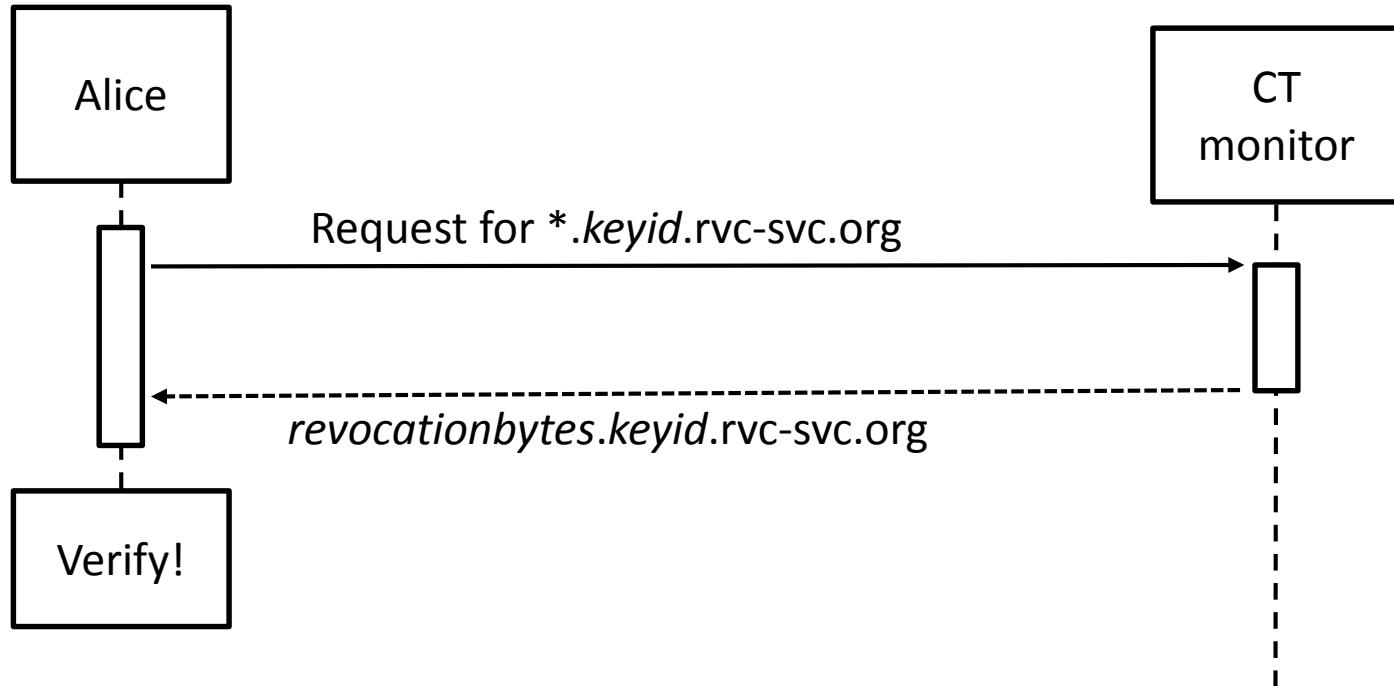
- Naively, clients store their revocation as specially encoded hostname
  - E.g. for [foo@bar.com](mailto:foo@bar.com) as *revocationbytes.foo.at.bar.com*
  - Then request a certificate for the hostname
- Clients can then inspect the log for those hostnames
- CT log is huge (~600MB)
- Optimise:
  - Provide well-known suffix, i.e. `rvc-svc.org`



# Revocation Publishing Protocol



# Revocation Querying Protocol



## Evaluation

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- + No PII
- + Same security as CT
  - o DNS label length
  - o CT logs cease to exist
- - searching the log leaks recipient



# Thanks!

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- Running SKS incurs legal and operational risks
- Users can be equivocated on or served malicious data
- Blockchain-based revocation scheme for OpenPGP certificates
  
- Contact
  - Tobias Mueller <mueller@informatik.uni-hamburg.de>
    - 0564 46F0 7732 1A69 1C67 14EA 8A01 4674 C937 42FD
  
  - Marius Stübs <stuebs@informatik.uni-hamburg.de>
    - 5887 4C3C CB21 3B1B F1A4 4A46 47D5 016C 42FF 7C2C